## 01204213: Homework 7

Due: 23pm, 20 Sep 2021.

1. (Sipser 5.1) In this problem we consider problems related to context free grammars. Let

 $ALL_{\mathsf{CFG}} = \{ \langle G \rangle \mid G \text{ is a CFG and } L(G) = \Sigma^*, \}$ 

be the language of all CFGs that accepts all strings. Also let

 $EQ_{\mathsf{CFG}} = \{ \langle G_1, G_2 \rangle \mid G_1 \text{ and } G_2 \text{ are CFGs and } L(G_1) = L(G_2), \}$ 

be the language of equivalent CFGs.

It is proved in the book that  $ALL_{CFG}$  is undecidable. Use this fact to prove that  $EQ_{CFG}$  is undecidable.

- 2. (Sipser 5.2) Show that  $EQ_{CFG}$  is co-Turing recognizable.
- 3. (Sipser 5.9) Let  $T = \{ \langle M \rangle \mid M \text{ is a TM that accepts } w^{\mathcal{R}} \text{ whenever it accepts } w \}$ . Show that T is undecidable.
- 4. (Sipser 5.23) Show that A is decidable iff  $A \leq_m 0^* 1^*$ .
- 5. (Sipser 5.22) Show that A is Turing-recognizable iff  $A \leq_m A_{\mathsf{TM}}$ .

Hint: Recall that a language A is Turing-recognizable iff there exists a TM M such that for every string w, M accepts w iff  $w \in A$ . You can use this definition in your solution. More specifically, you may want to use the fact that there exists a TM M in this definition.