Name, Binding, and Reference

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- Computer system manipulate objects.
 - Two ways one object can use another.
 - Use by value: Create a copy of the other object and pass that to the first object.
 - Use by reference: Name the other object and pass the name.
- Why use name?
 - Use by value does not permit **sharing.**
 - Allow designer to defer which object the name refers to until later.
 - Naming model
 - Operations
 - value <- RESOLVE(name, context)
 - status <- BIND(name, context)
 - list <- ENUMERATE(context)</p>
 - result <- COMPARE(name1, name2)</pre>
 - Vocabs
 - value = object or another name
 - *binding* = a map from a particular name to a particular value
 - name space = an alphabet of symbols together with syntac rules that specify which names are acceptable.
 - name-mapping algorithm = an algorithm that associates some names of the name space with some values
 - *resolve* = to perform the name-mapping algorithm from a name to the corresponding value
 - context = one of the inputs to the name-mapping algorithm. The common form is a set of name-value bindings
 - universal name space = a name space that has only one context. No matter who utters the name, the name has the same binding.
 - free name = a name that is not bound in a context
 - stable binding = a binding that stays the same for the lifetime of the namespace
- Three types of name mapping algorithms
 - Simple table lookup
 - Path name resolution
 - Search
- Default and explicit context reference
 - Who supplies the context argument in RESOLVE(name, context)?
 - *Default context reference* = the resolve supplies the context reference
 - Constant built into the resolver
 - Variable from the current environment
 - Explicit = specified by the object
 - Per object
 - Per name (*quantified name*)
 - Actually, the context argument is not the object. It is itself is a name, called the *context reference*. Resolver must resolve the context reference to a real context object. So, resolving is recursive. This recursion must end somewhere.
 - Problems
 - Object that utters the name does not provide an explicit context, and the name resolver chooses the wrong context.
 - Different contexts might bind different names for the same object.
 - □ Hard to pass name from one context to another.
 - □ Example: Phone numbers
- Path names and naming networks
 - Path name

- Explicitly include a reference to the context in which it should be resolved.
- Have multiple components.
 - □ Least significant component = the name
 - \Box All other parts = reference to the context in which the name is to be resolved.
- Recursive! Need to resolve the reference to the context (which itself is a path name) before we can resolve the name.
- Default context references for path name (the recursion must end somewhere!)
 - *Root* context built in to the resolver.
 - The resolver can store the path name of the default context. This needs to be resolved again. (Think working directory).
- *Absolute path name* = resolved using root context.
- *Relative path name* = resolved using the path name to the default context.
- Naming network
 - Contexts are objects.
 - Each context may bind a name to other contexts.
 - Name resolve chooses one context as root, and traces a path from the root to the first named context in the path name, and then the next, and so on until the path name runs out.
 - Different objects can have different path names in a naming network.
 - Names that refer to the same objects are synonyms.
 - Users may express path names relative to different roots. So it's hard for them to share names.
- Naming hierarchy
 - Naming network that is actually a tree.
 - Root context is the root of the tree.
 - Every object has a unique path name.
 - Very constraining. Not found in practice.
- Can add *indirect names* (names that are resolved to another path name) to naming hierarchy to permit cross-hierarchy linking.
- Search
 - \circ $\;$ Use an ordered list of contexts instead of a single default context.
 - Name resolver tries to resolve the name in the first context. If the result is not-found, then it tries the next context. This repeats until the list is exhausted or the resolver finds the first context that the name has a binding.
 - Search path is usually implemented as a per user list. This permits *user-dependent binding*. (Think about PATH variable.)
- Context Layers
 - Found in programming languages
 - Context are arranged into layers.
 - When the resolver cannot resolve a name in some layer, it tries resolving the name in the enclosing layer.
 - *Scope* = the range of layers a name is bound to the same object.
 - *Global name* = a name that is bound only in the outermost layer.
- Name discovery
 - How did you know to use this name?
 - *Name discovery protocol* = inform an object's importer the name that the object exports.
 - Exporter *advertise* the existence of the name.
 - Importer searches for an appropriate advertisement.
 - Forms
 - Well-known name = a name advertised so widely that it can be counted on to be stable.
 - Broadcast
 - Search = ask google
 - Reverse broadcast = ask everyone whether he/she knows the name of something
 - Narrowcast: Send "hello my name is ..." down the wire and hope that the other end will listen and reply.
 - Introduction = party and dating services

- Physical rendezvous
- Names and modular sharing
 - Modular sharing = can use an object, which itself is modular, without knowing the names of modules it uses.
 - Name conflicts is a syndrome of the lack of modular sharing.
 - Serious problem because resolving it means changing ways in which objects use name. This means changing the object itself!
 - *Imposed names* = names chosen by someone else
 - Common way to provide modular sharing: give each object its own context, and figure out a way to cross reference between contexts.
 - Programming languages use static scoping and closure to solve this problem. This mechanism is not found in file systems.
- Metadata and name overloading
 - *Metadata* = information about the object that is not a part of the object itself.
 - Name
 - Location
 - Time modified
 - Etc.
 - Overloaded name = name that has metadata in it
 - hello.c
 - Leonardo da Vinci
 - Physical address
 - ZIP Code
 - *Pure name* = no metadata inside, so has no relation to the object it refers to.
 - Problems with overloaded name.
 - Unstable name
 - Directory name /disk05. What if you move all the files to Disk 4 instead?
 - Tension between name stability and the need to update the overloaded information.
 - Address
 - Name of a physical location or a virtual location that maps to physical location.
 - Addresses are always overloaded. It always has information about the physical location of the object being referred to.
 - Address adjacency and physical adjacency go together. Arithmetic on addresses have corresponding physical meanings.
 - Addresses are extremely unstable.
 - How to cope with unstable names?

Hide unstable name with indirection.

- Have user of the object refer to the object by a generic name.
- The generic name itself is bound to the unstable name.
- Can change the generic-unstable-name binding later.
- Every problem in computer science can be solved with another layer of indirection.
- Generating unique names

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- Use consecutive integers as names.
 - Choose names at random from a large name space.
 - It's hard for a real machine to be perfectly random.
- Use hash functions.
 - But names generated this way is very unstable.
- Hierarchical naming scheme.
 - Useful for assigning names in a geographically distributed system.
 - Exploit delegation
 - Examples
 - Host names.
 - MAC addresses
- Relative lifetime of names, values, and bindings
 - Dangling name

- A name that outlives its binding.
- Resolve to irrelevant values or not-found.
- Can lead to serious errors when names are reused.
- Can be dealt with by verifying the object resulted from resolution if it meets the name user's expectation.

• Orphan

- An object that outlives its binding.
- Cannot be accessed by names again.
- Can be dealt with by reference counting or garbage collection.