## Name, Binding, and Reference

Monday, November 26, 2007
9:16 PM

- Computer system manipulate objects.
- Two ways one object can use another.
- Use by value: Create a copy of the other object and pass that to the first object.
- Use by reference: Name the other object and pass the name.
- Why use name?
- Use by value does not permit sharing.
- Allow designer to defer which object the name refers to until later.
- Naming model
- Operations
- value <- RESOLVE(name, context)
- status <- BIND(name, context)
- list <- ENUMERATE(context)
- result <- COMPARE(name1, name2)
- Vocabs
- value = object or another name
- binding = a map from a particular name to a particular value
- name space = an alphabet of symbols together with syntac rules that specify which names are acceptable.
- name-mapping algorithm = an algorithm that associates some names of the name space with some values
- resolve $=$ to perform the name-mapping algorithm from a name to the corresponding value
- context = one of the inputs to the name-mapping algorithm. The common form is a set of name-value bindings
- universal name space = a name space that has only one context. No matter who utters the name, the name has the same binding.
- free name $=$ a name that is not bound in a context
- stable binding $=$ a binding that stays the same for the lifetime of the namespace
- Three types of name mapping algorithms
- Simple table lookup
- Path name resolution
- Search
- Default and explicit context reference
- Who supplies the context argument in RESOLVE(name, context)?
- Default context reference $=$ the resolve supplies the context reference
- Constant built into the resolver
- Variable from the current environment
- Explicit = specified by the object
- Per object
- Per name (quantified name)
- Actually, the context argument is not the object. It is itself is a name, called the context reference. Resolver must resolve the context reference to a real context object. So, resolving is recursive. This recursion must end somewhere.
- Problems
- Object that utters the name does not provide an explicit context, and the name resolver chooses the wrong context.
- Different contexts might bind different names for the same object.
$\square$ Hard to pass name from one context to another.
- Example: Phone numbers
- Path names and naming networks
- Path name
- Explicitly include a reference to the context in which it should be resolved.
- Have multiple components.
- Least significant component $=$ the name
- All other parts = reference to the context in which the name is to be resolved.
- Recursive! Need to resolve the reference to the context (which itself is a path name) before we can resolve the name.
- Default context references for path name (the recursion must end somewhere!)
- Root context built in to the resolver.
- The resolver can store the path name of the default context. This needs to be resolved again. (Think working directory).
- Absolute path name $=$ resolved using root context.
- Relative path name $=$ resolved using the path name to the default context.
- Naming network
- Contexts are objects.
- Each context may bind a name to other contexts.
- Name resolve chooses one context as root, and traces a path from the root to the first named context in the path name, and then the next, and so on until the path name runs out.
- Different objects can have different path names in a naming network.
- Names that refer to the same objects are synonyms.
- Users may express path names relative to different roots. So it's hard for them to share names.
- Naming hierarchy
- Naming network that is actually a tree.
- Root context is the root of the tree.
- Every object has a unique path name.
- Very constraining. Not found in practice.
- Can add indirect names (names that are resolved to another path name) to naming hierarchy to permit cross-hierarchy linking.
- Search
- Use an ordered list of contexts instead of a single default context.
- Name resolver tries to resolve the name in the first context. If the result is not-found, then it tries the next context. This repeats until the list is exhausted or the resolver finds the first context that the name has a binding.
- Search path is usually implemented as a per user list. This permits user-dependent binding. (Think about PATH variable.)
- Context Layers
- Found in programming languages
- Context are arranged into layers.
- When the resolver cannot resolve a name in some layer, it tries resolving the name in the enclosing layer.
- Scope $=$ the range of layers a name is bound to the same object.
- Global name = a name that is bound only in the outermost layer.
- Name discovery
- How did you know to use this name?
- Name discovery protocol = inform an object's importer the name that the object exports.
- Exporter advertise the existence of the name.
- Importer searches for an appropriate advertisement.
- Forms
- Well-known name = a name advertised so widely that it can be counted on to be stable.
- Broadcast
- Search = ask google
- Reverse broadcast = ask everyone whether he/she knows the name of something
- Narrowcast: Send "hello my name is ..." down the wire and hope that the other end will listen and reply.
- Introduction = party and dating services
- Physical rendezvous
- Names and modular sharing
- Modular sharing = can use an object, which itself is modular, without knowing the names of modules it uses.
- Name conflicts is a syndrome of the lack of modular sharing.
- Serious problem because resolving it means changing ways in which objects use name. This means changing the object itself!
- Imposed names = names chosen by someone else
- Common way to provide modular sharing: give each object its own context, and figure out a way to cross reference between contexts.
- Programming languages use static scoping and closure to solve this problem. This mechanism is not found in file systems.
- Metadata and name overloading
- Metadata $=$ information about the object that is not a part of the object itself.
- Name
- Location
- Time modified
- Etc.
- Overloaded name $=$ name that has metadata in it
- hello.c
- Leonardo da Vinci
- Physical address
- ZIP Code
- Pure name = no metadata inside, so has no relation to the object it refers to.
- Problems with overloaded name.
- Unstable name
- Directory name /disk05. What if you move all the files to Disk 4 instead?
- Tension between name stability and the need to update the overloaded information.
- Address
- Name of a physical location or a virtual location that maps to physical location.
- Addresses are always overloaded. It always has information about the physical location of the object being referred to.
- Address adjacency and physical adjacency go together. Arithmetic on addresses have corresponding physical meanings.
- Addresses are extremely unstable.
- How to cope with unstable names?
- Hide unstable name with indirection.
- Have user of the object refer to the object by a generic name.
- The generic name itself is bound to the unstable name.
- Can change the generic-unstable-name binding later.
- Every problem in computer science can be solved with another layer of indirection.
- Generating unique names
- Use consecutive integers as names.
- Choose names at random from a large name space.
- It's hard for a real machine to be perfectly random.
- Use hash functions.
- But names generated this way is very unstable.
- Hierarchical naming scheme.
- Useful for assigning names in a geographically distributed system.
- Exploit delegation
- Examples
- Host names.
- MAC addresses
- A name that outlives its binding.
- Resolve to irrelevant values or not-found.
- Can lead to serious errors when names are reused.
- Can be dealt with by verifying the object resulted from resolution if it meets the name user's expectation.
- An object that outlives its binding.
- Cannot be accessed by names again.
- Can be dealt with by reference counting or garbage collection.

